

CS7800: Advanced Algorithms

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Linear Programming (Part 2)

Duality

Recall: The Brewery LP

objective function

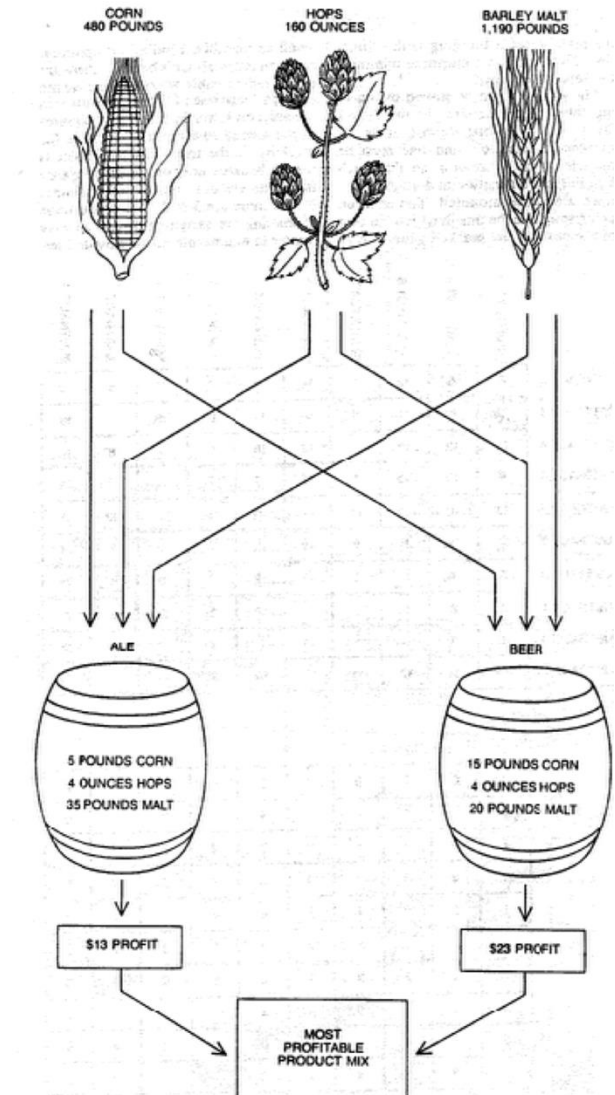
Ale Beer

$$\begin{array}{ll} \max & 13A + 23B \\ \text{s. t.} & 5A + 15B \leq 480 \\ & 4A + 4B \leq 160 \\ & 35A + 20B \leq 1190 \\ & A, B \geq 0 \end{array}$$

constraint

decision variable

Beverage	Corn (pounds)	Hops (ounces)	Malt (pounds)	Profit (\$)
Ale (barrel)	5	4	35	13
Beer (barrel)	15	4	20	23
constraint	480	160	1190	



Duality

“Primal” problem:

$$\begin{array}{llllll} \text{(P)} & \max & 13A & + & 23B & \\ & \text{s. t.} & 5A & + & 15B & \leq 480 \\ & & 4A & + & 4B & \leq 160 \\ & & 35A & + & 20B & \leq 1190 \\ & & A & , & B & \geq 0 \end{array}$$

Goal 1: Find a lower bound on optimal value.

Easy: Any feasible solution provides one.

- $(A, B) = (34, 0) \Rightarrow z^* \geq 442$
- $(A, B) = (0, 32) \Rightarrow z^* \geq 736$
- $(A, B) = (7.5, 29.5) \Rightarrow z^* \geq 776$
- $(A, B) = (12, 28) \Rightarrow z^* \geq 800$

Duality

“Primal” problem:

$$\begin{array}{llllll} \text{(P)} & \max & 13A & + & 23B & \\ & \text{s. t.} & 5A & + & 15B & \leq 480 \\ & & 4A & + & 4B & \leq 160 \\ & & 35A & + & 20B & \leq 1190 \\ & & A & , & B & \geq 0 \end{array}$$

Goal 2: Find an **upper bound** on optimal value.

Ex 1: Multiply 2nd inequality by 6: $24A + 24B \leq 960$. This gives

$$z^* = 13A + 23B \leq 24A + 24B \leq 960.$$

Ex 2: Add two times 1st inequality to 2nd inequality:

$$z^* = 13A + 23B \leq 14A + 34B \leq 1120.$$

Duality

“Primal” problem:

$$\begin{array}{llllll} \text{(P)} & \max & 13A & + & 23B & \\ & \text{s. t.} & 5A & + & 15B & \leq 480 \\ & & 4A & + & 4B & \leq 160 \\ & & 35A & + & 20B & \leq 1190 \\ & & A & , & B & \geq 0 \end{array}$$

Goal 2: Find an upper bound on optimal value.

Ex 3: Add one time 1st inequality to two times 2nd inequality:

$$z^* = 13A + 23B \leq 13A + 23B \leq 800.$$

Recall lower bound: $(A, B) = (12, 28) \Rightarrow z^* \geq 800$

Combining upper and lower bounds: $z^* = 800$.

Duality

$$\begin{array}{llllll} \text{(P)} & \max & 13A & + & 23B & \\ & \text{s. t.} & 5A & + & 15B & \leq 480 \\ & & 4A & + & 4B & \leq 160 \\ & & 35A & + & 20B & \leq 1190 \\ & & A & , & B & \geq 0 \end{array}$$

- Let's try to generalize this idea...
- **Idea:** Add non-negative combination (C, H, M) of constraints s.t.

$$\begin{aligned} 13A + 23B &\leq (5C + 4H + 35M) A + (15C + 4H + 20M) B \\ &\leq 480C + 160H + 1190M \end{aligned}$$

- But how do we find C, H, M ?

Duality

$$\begin{array}{llllll} \text{(P)} & \max & 13A & + & 23B & \\ & \text{s. t.} & 5A & + & 15B & \leq 480 \\ & & 4A & + & 4B & \leq 160 \\ & & 35A & + & 20B & \leq 1190 \\ & & A & , & B & \geq 0 \end{array}$$

- Let's try to generalize this idea...
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- But how do we find C, H, M ? By writing another LP... 🤖
- Dual:** Find the best upper bound with constraint coefficients.

$$\begin{array}{llllll} \text{(D)} & \min & 480C & + & 160H & + & 1190M \\ & \text{s. t.} & 5C & + & 4H & + & 35M \geq 13 \\ & & 15C & + & 4H & + & 20M \geq 23 \\ & & C & , & H & , & M \geq 0 \end{array}$$

Duality: Economic Interpretation

- **Brewer / primal:** find optimal mix of beer and ale to maximize profit.

$$\begin{array}{llllll} \text{(P)} & \max & 13A & + & 23B & \\ & \text{s. t.} & 5A & + & 15B & \leq 480 \\ & & 4A & + & 4B & \leq 160 \\ & & 35A & + & 20B & \leq 1190 \\ & & A & , & B & \geq 0 \end{array}$$

- **Entrepreneur / dual:** buy each resource from brewer at min cost.
 - C, H, M : unit price for corn, hops, malt.
 - Brewer won't sell if $5C + 4H + 3M < 13$.

$$\begin{array}{llllll} \text{(D)} & \min & 480C & + & 160H & + & 1190M \\ & \text{s. t.} & 5C & + & 4H & + & 35M \geq 13 \\ & & 15C & + & 4H & + & 20M \geq 23 \\ & & C & , & H & , & M \geq 0 \end{array}$$

Duality: Basic Observations

- If the primal has n variables and m constraints, then the dual has m variables and n constraints.

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$$\begin{array}{llllllll} \text{(D)} & \min & 480C & + & 160H & + & 1190M & \\ & \text{s. t.} & 5C & + & 4H & + & 35M & \geq 13 \\ & & 15C & + & 4H & + & 20M & \geq 23 \\ & & C & , & H & , & M & \geq 0 \end{array}$$

Double dual

- Canonical form:

$$\begin{array}{ll} \text{(P)} & \max c^T x \\ & \text{s. t. } Ax \leq b \\ & x \geq 0 \end{array}$$

$$\begin{array}{ll} \text{(D)} & \min y^T b \\ & \text{s. t. } A^T y \geq c \\ & y \geq 0 \end{array}$$

- **Property:** The dual of the dual is the primal.
- **Pf:** Rewrite (D) as maximization in canonical form, take its dual.

$$\begin{array}{ll} \text{(D')} & \max -y^T b \\ & \text{s. t. } -A^T y \leq -c \\ & y \geq 0 \end{array}$$

$$\begin{array}{ll} \text{(DD)} & \min -c^T z \\ & \text{s. t. } -(A^T)^T z \geq -b \\ & z \geq 0 \end{array}$$

The Dual Recipe

Primal (P)	maximize
constraints	$a x = b_i$ $a x \leq b$ $a x \geq b_i$
variables	$x_j \geq 0$ $x_j \leq 0$ unrestricted

minimize	Dual (D)
y_i unrestricted $y_i \geq 0$ $y_i \leq 0$	variables
$a^T y \geq c_j$ $a^T y \leq c_j$ $a^T y = c_j$	constraints

- Pf: Rewrite in canonical form and take dual.

Strong Duality

- **Theorem (Weak Duality):** For any A, b, c if (P) and (D) are feasible, then $\max \leq \min$.

$$\begin{array}{ll} \text{(P)} & \max c^T x \\ & \text{s. t. } Ax = b \\ & \quad x \geq 0 \end{array}$$

$$\begin{array}{ll} \text{(D)} & \min y^T b \\ & \text{s. t. } A^T y \geq c \end{array}$$

- **Theorem (Strong Duality):** For any A, b, c if (P) and (D) are feasible, then $\max = \min$.

Examples: Single Source Shortest Path

- **Input:** Weighted directed graph $G=(V, E)$, weight function $\ell: E \rightarrow R$ and vertex $s \in V$.
- **Output:** Length of the shortest path from s to every other vertex.

$$\begin{array}{ll} \text{maximize} & \sum_v \text{dist}(v) \\ \text{subject to} & \text{dist}(s) = 0 \\ & \text{dist}(v) - \text{dist}(u) \leq \ell(u \rightarrow v) \quad \text{for every edge } u \rightarrow v \end{array}$$

- What's the dual?